

**USC GridSec Project**  
**Trusted Grid Computing with**  
**Distributed Security Control**  
**and Defense against Intrusions**

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**Research Meeting on February 10, 2004**

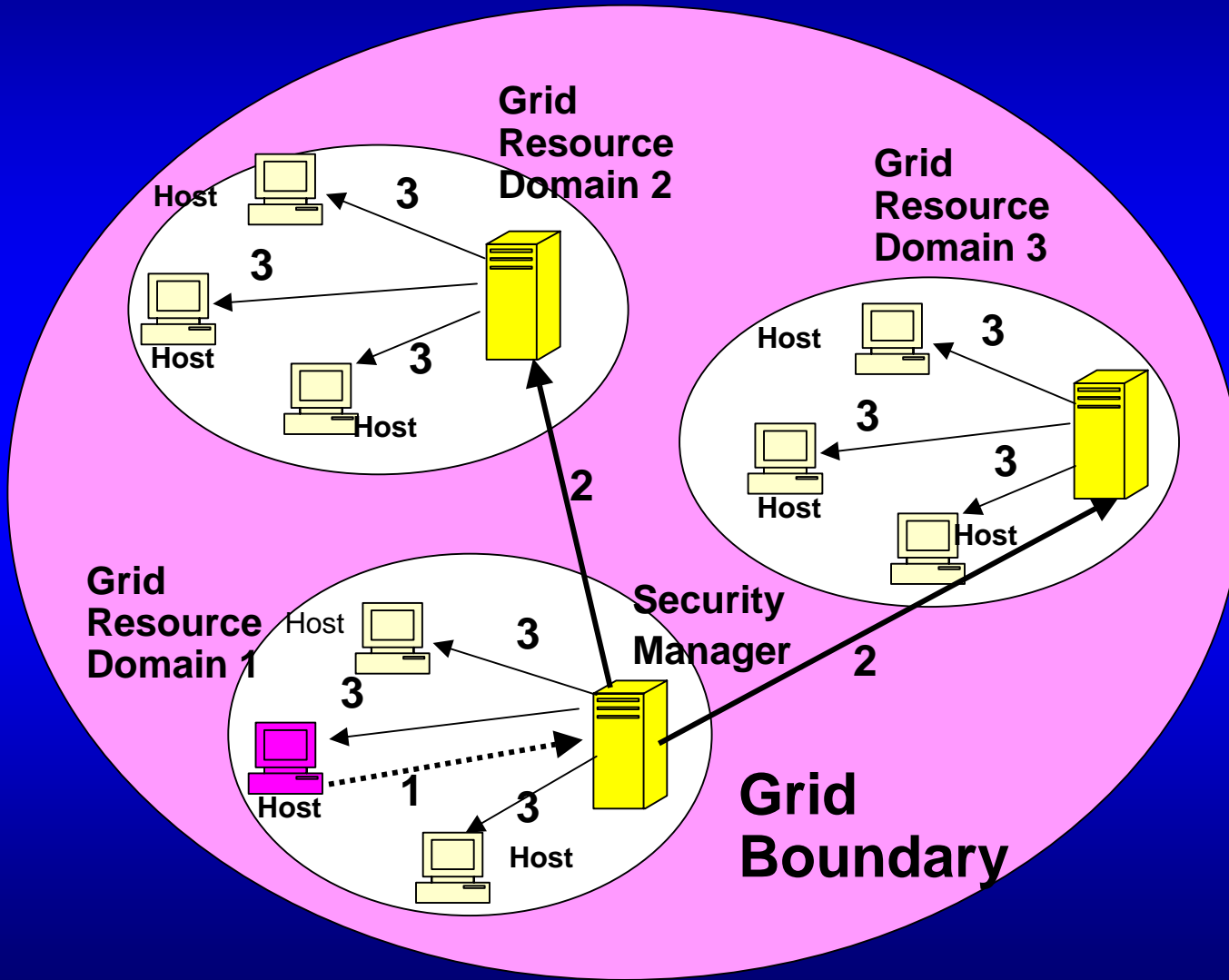
# Meeting Agenda :

- 1. The GridSec Project at USC**
  - Hwang and Team
- 2. Grid Resource Access-Control and Policy Management**
  - Neuman and Ryutov
- 3. FPGA-based Network Security Support**
  - Prasanna and Baker
- 4. Discussions on Planned Future Work**

# The GridSec Project at USC:

- Sponsored by a NSF/ITR Research Grant ACR-0325409
- PI and Project Director: Kai Hwang  
Co-PIs: Clifford Neuman and Viktor Prasanna
- Post-doctorial and Foreign Collaborators:  
Dr. Tatyana Ryutov of USC/ISI  
Prof. Michel Cosnard of INRIA, France  
Dr. Zhiwei Xu of Chinese Academy of Sciences
- Ph.D.-bound Research Assistants:  
M. Qin, S. Song, Y. Kim, Z. Baker, L. Zhou,  
Y. Zhang, R. Tripathi, C. Chuan, .....

# Distributed GridSec Architecture



**Step 1:** .....➔  
Intrusion detected by a local micro-firewall

**Step 2:** ➔➔  
All security managers alerted with the intrusion

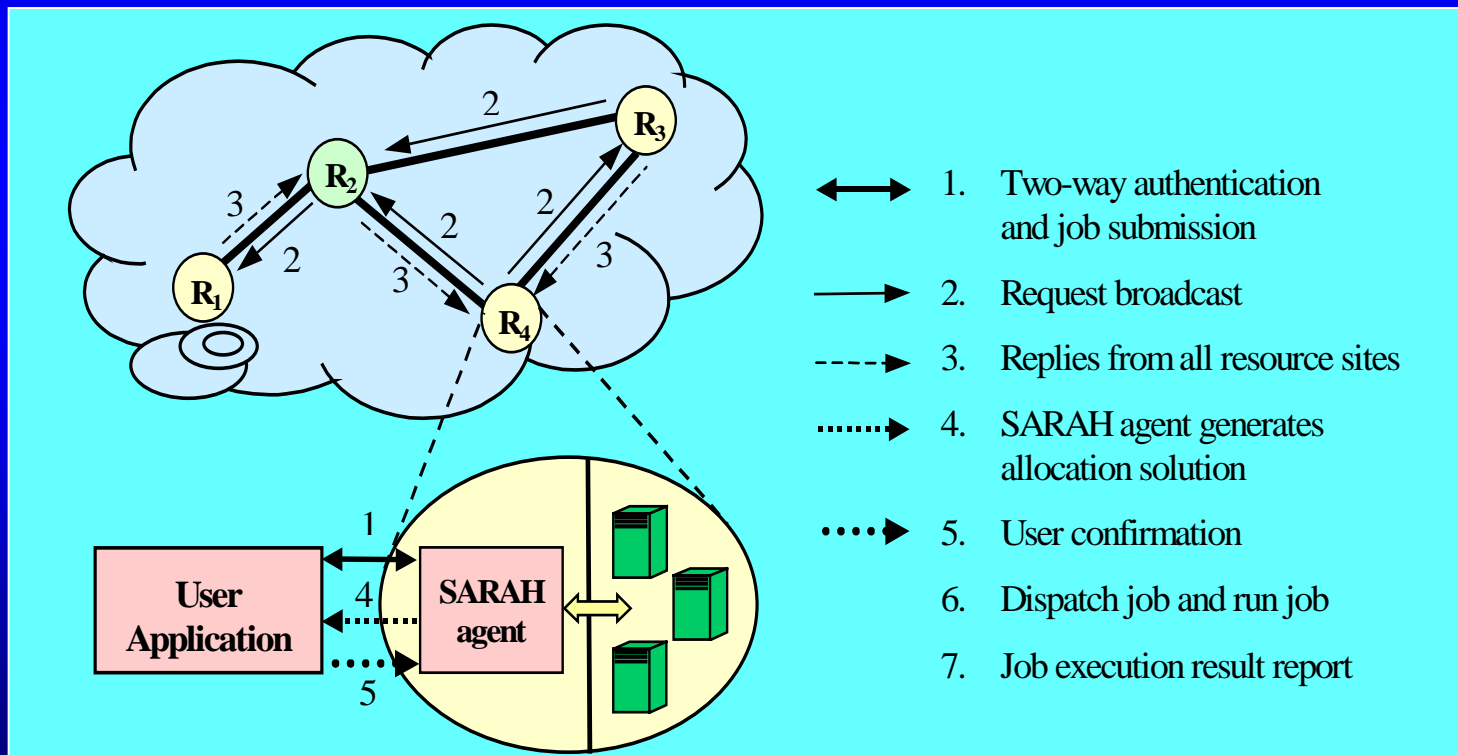
**Step 3:** ➔➔➔  
Security managers broadcast response command to all hosts under their jurisdiction.

## Building Encrypted Tunnels between Resource Networks or Clusters for Trusted Grid Computing

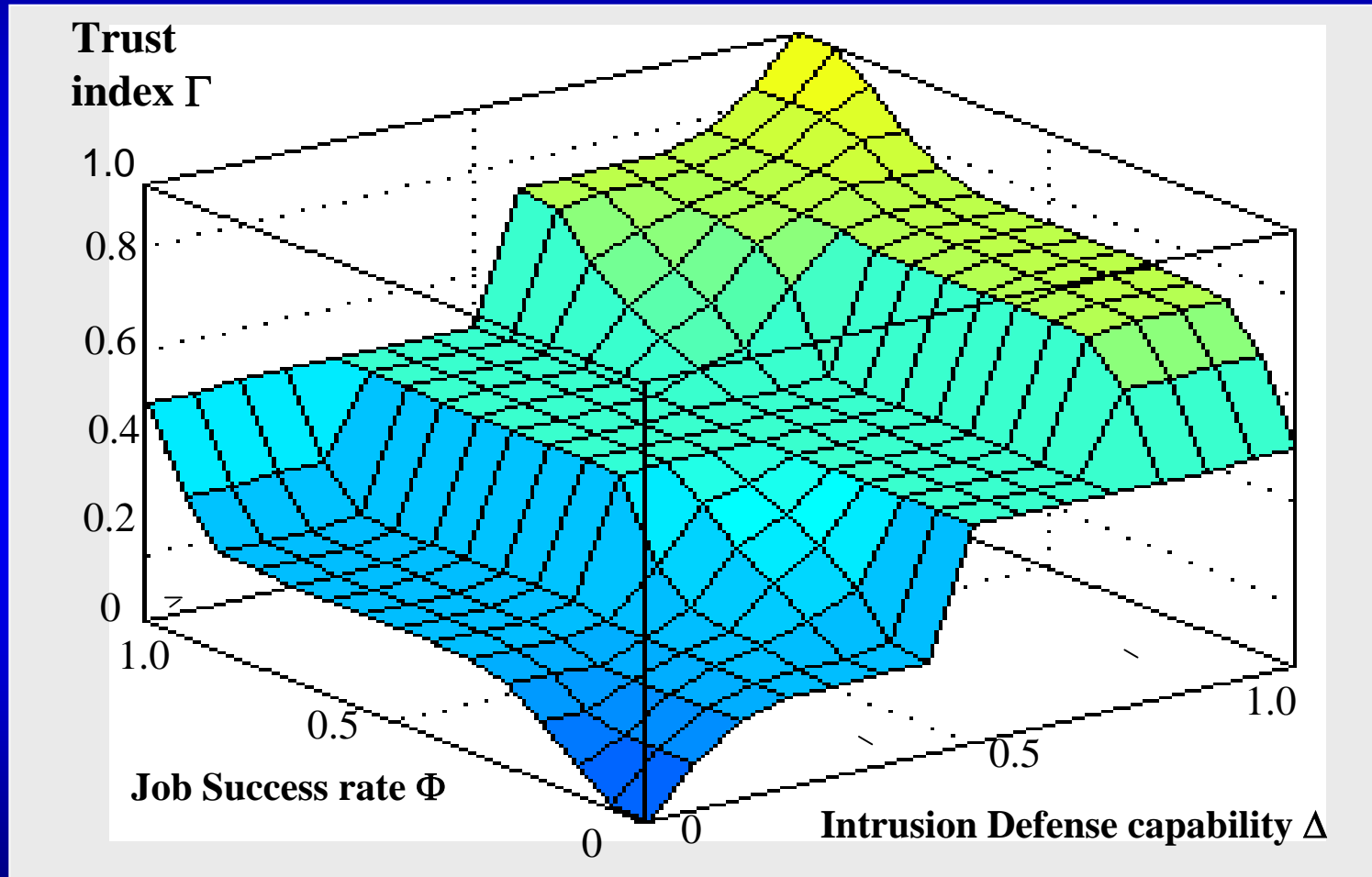
- With  $n$  nodes, the required number of tunnels grows with  $O(n^2)$
- By using Dijkstra algorithm (All-source shortest path ) in building VPN with minimum number of tunnels which grows with  $O(n \log n)$  complexity (Yongjin Kim)
- Upon shortest path, security policy is generated to satisfy special Grid requirements, automatically if possible.
- How to integrate security policies from various private networks through the public network ? (Gurpreet Grewal)
- How to resolve policy conflicts among hosts, firewalls, switches, routers, and servers, etc. in a Grid environment ?

# Security-Assured Resource Allocation (Shanshan Song and Mikin Macwan)

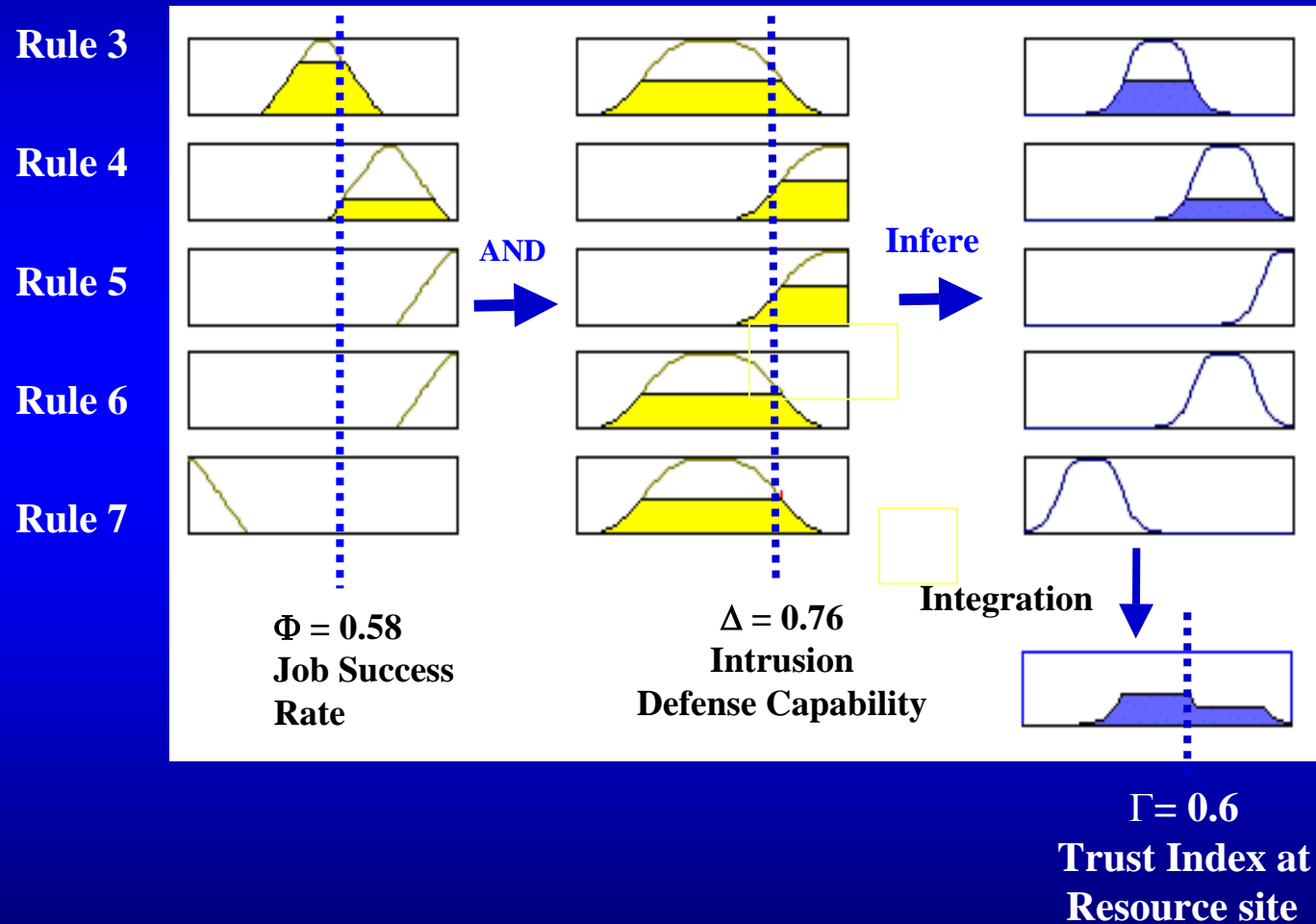
- Fuzzy logic for trust management over multiple Grid resource sites
- Linear programming model for trusted resource allocation (Sarah)



# Trust Index of Grid Resource Site: derived from job success rate and intrusion defense capability at resource site periodically



# Fuzzy Logic Inference for Assessing Trust Index at Grid Resource Site





# Security-Assured Resource Allocation

**Algorithm: Dual-objective linear programming for secure resource allocation**

**Inputs:** (1) A Grid job  $J = (P_o, T_o, C_o)$ , submitted to the  $j$ -th resource site  $R_j$   
(2)  $P_i$  and  $C_i$  ( $i = 1, 2, \dots, k$ ) for  $k$  resource sites and trust vector  $V_j$  in  $R_j$ .

**Output:** The resource allocation vector  $X = (x_1, x_2, \dots, x_k)$ .

**Objective Function:** Maximize  $P = \sum_{i=1}^k x_i P_i t_{ij}$  and minimize  $C = \sum_{i=1}^k x_i P_i C_i$  simultaneously.

or equivalently, maximize the performance/cost ratio  $E$  defined in Eq. (6).

**Subject to** the following constraints:

$$\sum_{i=1}^k x_i P_i \geq P_o, \quad \sum_{i=1}^k x_i P_i C_i \leq C_o, \quad \text{and} \quad t_{ij} \geq T_o \quad (7)$$

where  $0 \leq x_i \leq 1$  for  $i = 1, 2, \dots, k$ .

# A SARAH Resource Optimization Example

$V_1$	$V_2$	$V_3$	$V_4$
0.5	0.5	0.4	0.4
0.8	0.8	0.7	1.0
0.9	1.0	0.8	0.8
0.7	0.7	0.6	0.6

(a) Trust Matrix

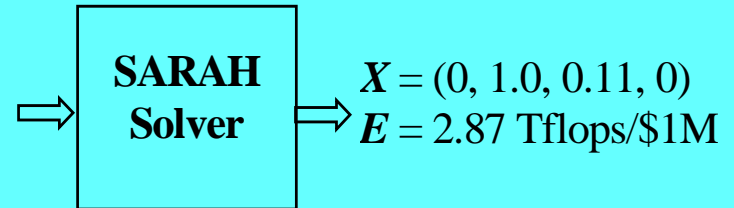
$$J_1 = (1.5\text{Tflops}, \$600\text{K}, 0.6)$$

$$P_1 = 1.1\text{Tflops}, C_1 = \$300\text{K}$$

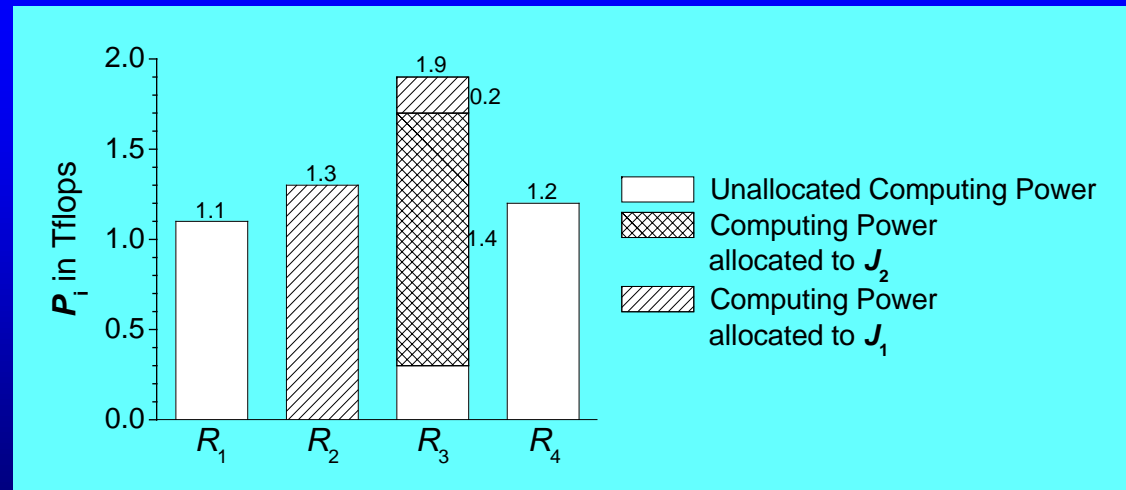
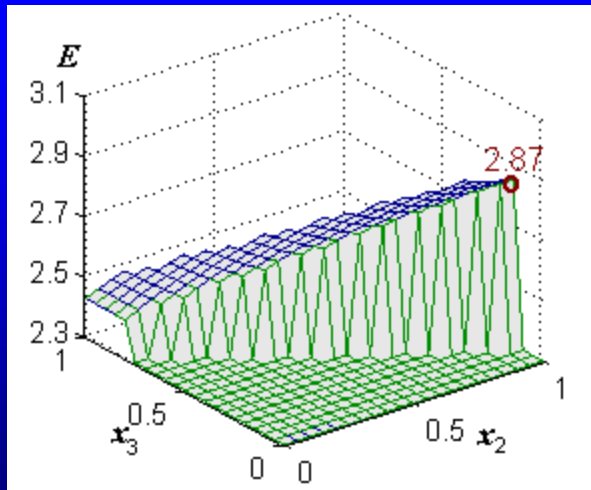
$$P_2 = 1.3\text{Tflops}, C_2 = \$340\text{K}$$

$$P_3 = 1.9\text{Tflops}, C_3 = \$330\text{K}$$

$$P_4 = 1.2\text{Tflops}, C_4 = \$300\text{K}$$



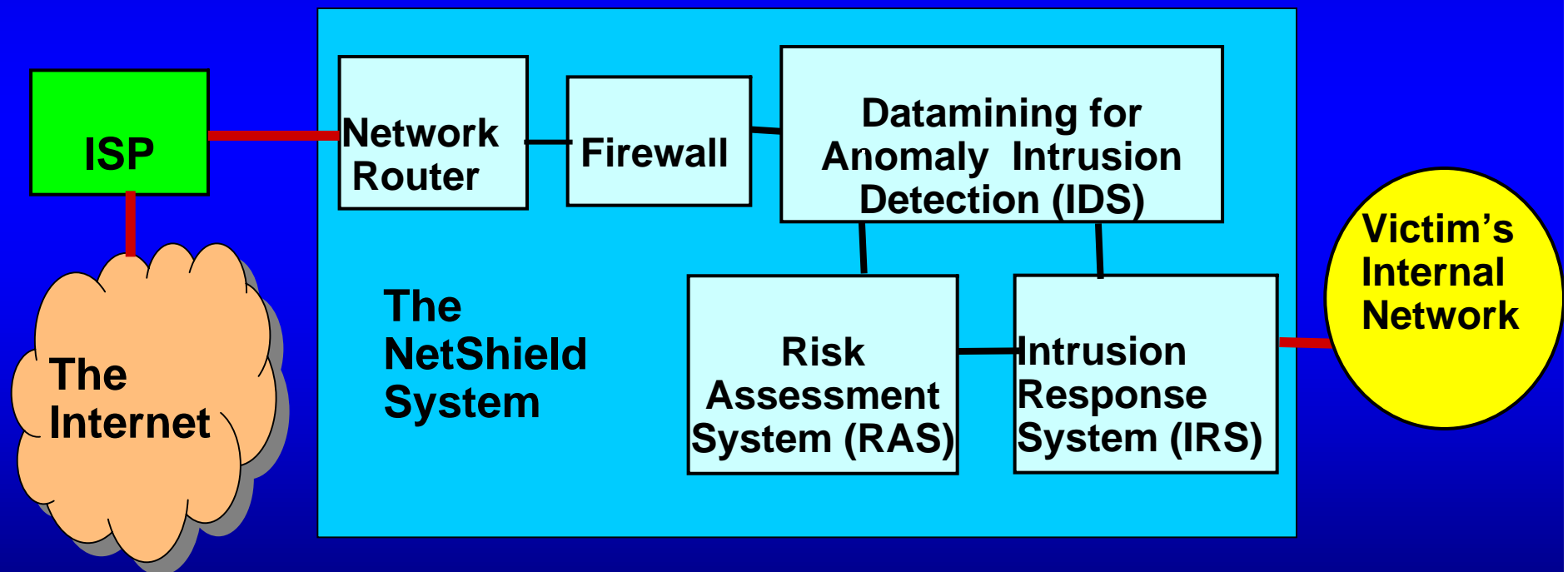
(b) Inputs and outputs



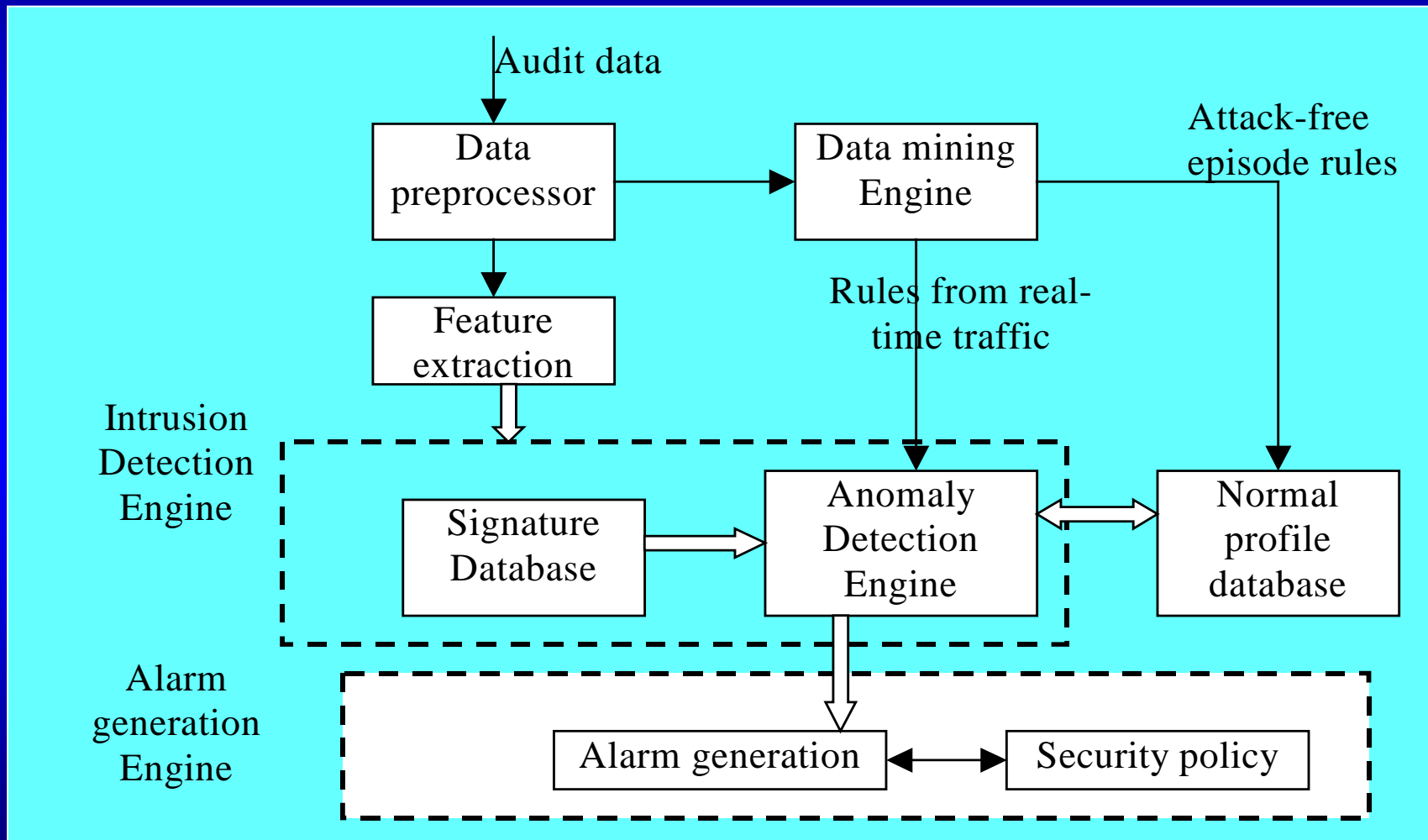
# USC NetShield Intrusion Defense System

## for Protecting Grid Computing Resource Sites

(Narayana, Chuan, and Qin)



# Data mining Architecture for Network Anomaly Detection (Min Qin)



# Datamining For Intrusion Detection

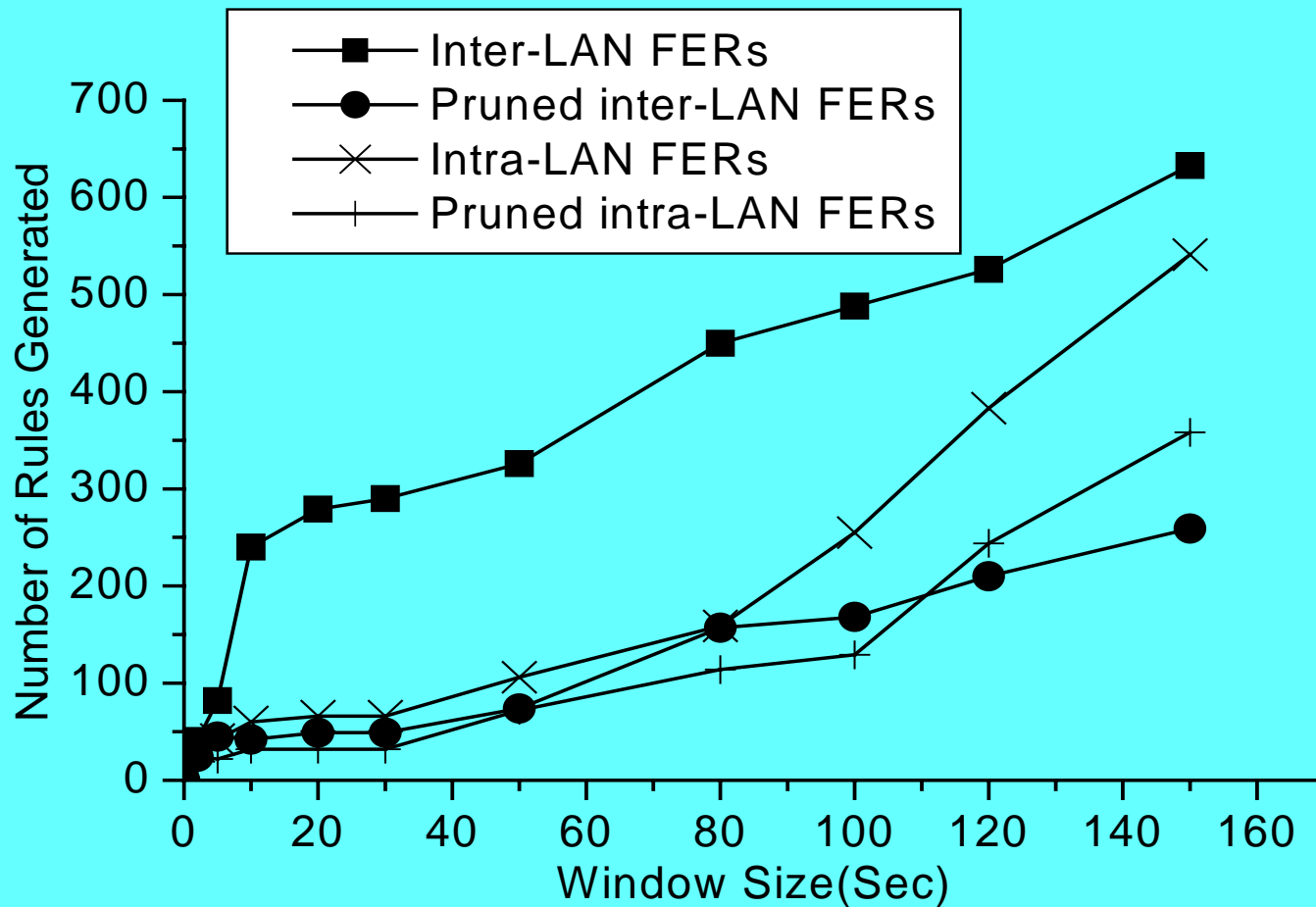
Trace TCP/UDP/ICMP connection events  
to derive frequent episode rule (FER)  
on traffic patterns. The SYN flood attack is  
specified by the following rule:

**(service = http, flag = S0) (service = http, flag = S0)**

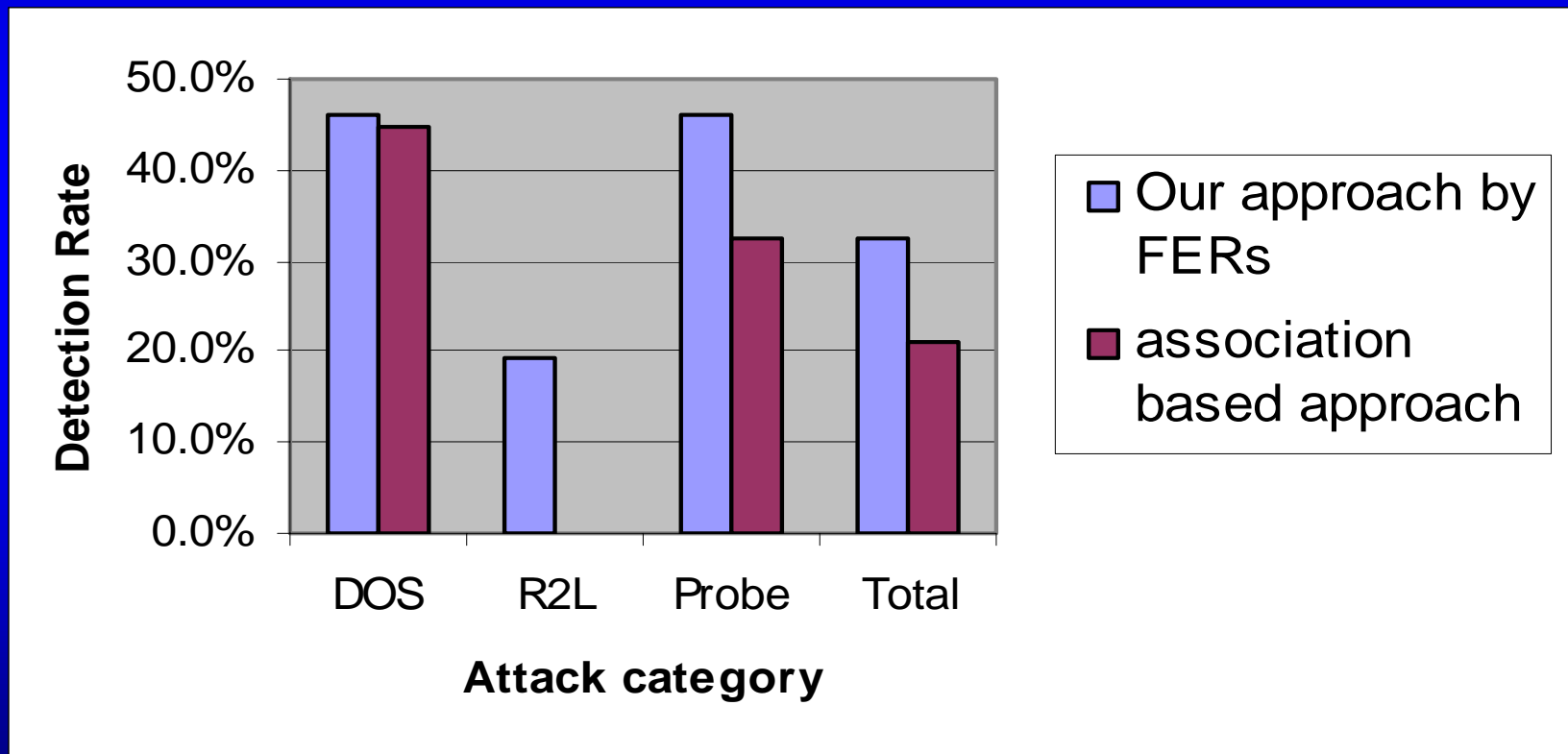
**→ (service = http, flag = S0)**

# Reduction of Rule Search Space

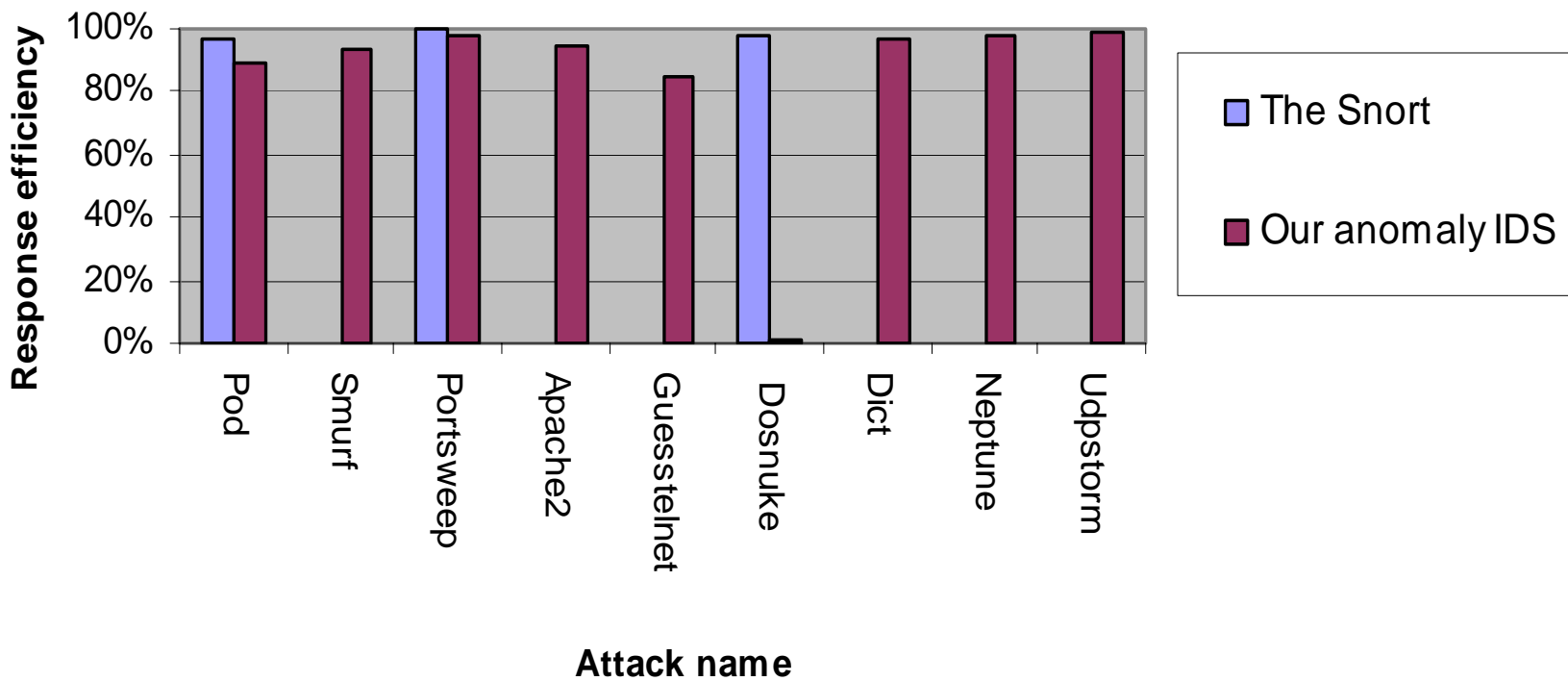
after pruning of ineffective FERs



# Anomaly Detection Results using Frequent Episode Rules



# Intrusion Response Efficiency





# DDoS Detection and Defense

(Y. Zhang and R. Tripathi)

- Internet traffic flow analysis for packet filtering, rate limiting, priority queuing, etc.
- Smart filtering for countering DDoS attacks based on traffic history analysis
- Use the Georgia Tech Network Simulator
- Test on the USC Internet trace data sets

## Recent Papers and Technical Reports :

1. M. Qin and K. Hwang, “ Frequent Episode Rules for Intrusive Anomaly Dtection with Internet Datamining”, *USENIX Security Symposium*, submitted Jan.27, 2004
2. S. Song and k. Hwang, “ Integrated Trust and Resource Optimization for Security-Assured Grip Computing”, *International Symposium on High-Performance Distributed Computing*, (HPDC-13), submitted Feb. 7, 2004
3. Z. Baker and V. Prasanna”, A Methodology for Synthesis of Efficient Intrusion Detection Systems on FPGAs”, *IEEE Field-Programmable Custom Comuting Mechanism (FCCM)*, submitted Jan. 19, 2004
4. K. Hwang, et al, “ GridSec and NetShield Architecture for Securing Grid Computing”, *Tech. Report 2003-1*, USC Internet and Grid Computing Lab., (under revision)